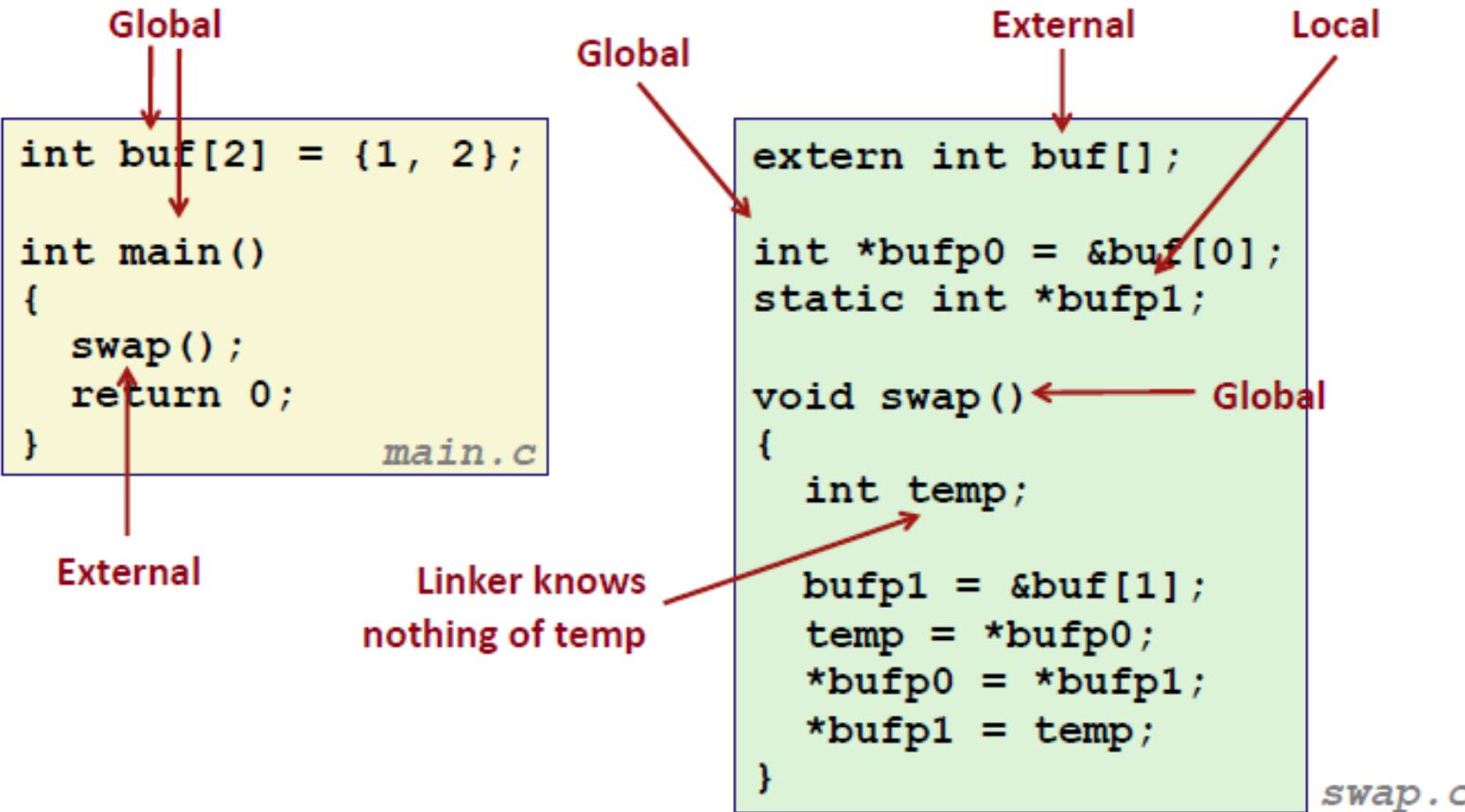


# Resolving Symbols



# Relocating Code and Data

main.c

```
int buf[2] = {1, 2};

int main()
{
    swap();
    return 0;
}
```

swap.c

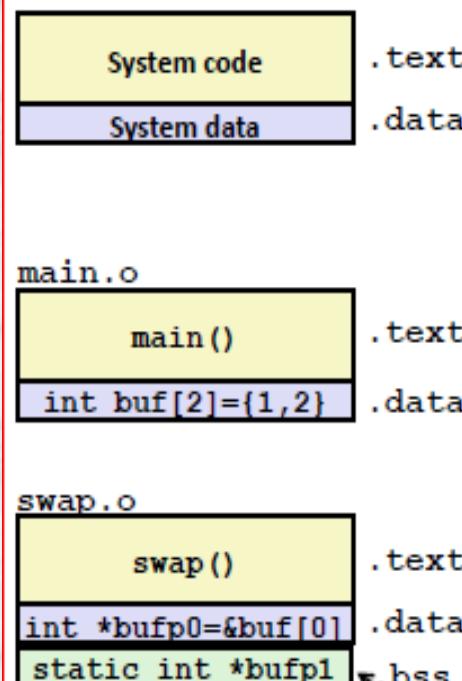
```
extern int buf[];

int *bufp0 = &buf[0];
static int *bufp1;

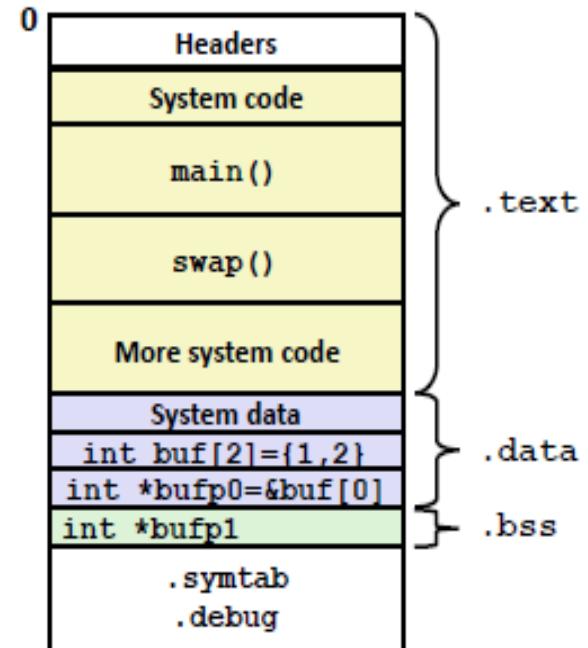
void swap()
{
    int temp;

    bufp1 = &buf[1];
    temp = *bufp0;
    *bufp0 = *bufp1;
    *bufp1 = temp;
}
```

Relocatable Object Files



Executable Object File



Even though private to swap, requires allocation in .bss

# Practice problem 7.1 (pg 662)

**swap.c**

```
extern int buf[];

int *bufp0 = &buf[0];
static int *bufp1;

void swap()
{
    int temp;

    bufp1 = &buf[1];
    temp = *bufp0;
    *bufp0 = *bufp1;
    *bufp1 = temp;
}
```

- SYMBOL TABLE → .symtab
  - Info about functions and global variables that are defined and referenced in a program
    - Does not contain entries for local variables

- Understanding the relationship between linker symbols and C variables/functions... notice that the C local variable temp does NOT have a symbol table entry. Why? It goes on the stack!

Symbol	swap.o .symtab entry?	symbol type	module where defined	section
buf	yes	extern	main.o	.data
bufp0	yes	global	swap.o	.data
bufp1	yes	global	swap.o	.bss
swap	yes	global	swap.o	.text
temp	no	---	---	---

# Swap relocatable symbol table

```
% objdump -r -d -t swap.o  
swap.o: file format elf32-i386  
  
SYMBOL TABLE:  
  
00000000 I df *ABS*      00000000 swap.c  
00000000 I d .text       00000000 .text  
00000000 I d .data       00000000 .data  
00000000 I d .bss        00000000 .bss  
00000000 I O .bss        00000004 bufp1  
00000000 g O .data        00000004 bufp0  
00000000 *UND*          00000000 buf  
00000000 g F .text        00000035 swap
```

O = object  
d = debug  
I = local

F = function  
f = file  
g = global

## swap.c

```
extern int buf[];  
  
int *bufp0 = &buf[0];  
static int *bufp1;  
  
void swap()  
{  
    int temp;  
  
    bufp1 = &buf[1];  
    temp = *bufp0;  
    *bufp0 = *bufp1;  
    *bufp1 = temp;  
}
```

# Symbols and Symbol Tables

- Local linker symbols != local program variables
  - ➊ .syms does not contain any symbols that correspond to local non-static program variables.
  - ➋ These are managed at run time on the stack and are not of interest to the linker
  - ➌ However... local procedure variables that are defined with the C static attribute (EXCEPTION) are not managed on the stack
    - The compiler allocates space in .data or .bss for each and creates a local linker symbol in the symbol table with a unique name

FYI: Static variable's lifetime extends across the entire run of the program where local variables are allocated and deallocated on the stack

# SYMBOL TABLES

- Built by assemblers using symbols exported by the compiler into the .s file
- An ELF symbol table is contained in the .syms table section
- It contains an array of entries where each entry contains:
  - Symbol's value i.e. address
  - Flag bits (l=local, g = global, F=function, etc)
    - Characters and spaces – up to 7 bits
  - Section or
    - \*ABS\* (absolute – not in any section)
    - \*UND\* if referenced but not defined
  - Alignment or size
  - Symbol's name

# Relocation

- Relocation merges the input modules and assigns run-time addresses to each symbol
- When an assembler generates an object module, it does not know where the code and data will ultimately be stored in memory or the locations of any externally defined functions or global variables referenced by the module
- A “relocation entry” is generated when the assembler encounters a reference to an object whose ultimate location is unknown
- 2 types
  - R\_386\_PC32
  - R\_386\_32

# 2 Relocation types

## R\_386\_PC32

-  relocate a reference that uses a 32-bit PC-relative address.
-  Effective address = PC + instruction encoded addr

## R\_386\_32

-  Absolute addressing
-  Uses the value encoded in the instruction

# Relocation Info

main.c

```
int buf[2] = {1, 2};

int main()
{
    swap();
    return 0;
}
```

```
% gcc -o main -m32 main.c
/tmp/ccEVrEUg.o:
In function `main':
main.c:(.text+0x7): undefined
reference to `swap'
collect2:
ld returned 1 exit status
```

Call instruction @ offset 0x6,  
opcode e8, 32-bit ref (-4)  
Offset = 0x7 Symbol = swap  
Type = R\_386\_PC32

```
% gcc -c -m32 main.c
% objdump -r -tdata main.o
```

SYMBOL TABLE:

00000000	I	df *ABS*	00000000	main.c
00000000	I	d .text	00000000	.text
00000000	I	d .data	00000000	.data
00000000	I	d .bss	00000000	.bss
00000000	g	O .data	00000008	buf
00000000	g	F .text	00000014	main
00000000		*UND*	00000000	swap

Disassembly of section .text: **0x80483b4**

00000000 <main>:

0:	55	push %ebp
1:	89 e5	mov %esp,%ebp
3:	83 e4 f0	and \$0xffffffff0,%esp
6:	e8 fc ff ff ff	call 7 <main+0x7>

**7: R\_386\_PC32 swap**

b:	b8 00 00 00 00	mov \$0x0,%eax
10:	89 ec	mov %ebp,%esp
12:	5d	pop %ebp
13:	c3	ret

.symbol swap + 32-bit ref – ref addr  
0x80483c8 + -4 - 0x80483bb = 0x9

NOTE: relocation entries and instructions are stored in different sections of the object file → .rel.txt vs .text

.text + .offset  
0x80483b4 + 0x7  
= 0x80483bb = ref addr

**0x80483c8**

call swap

Relocation entry

# Executable before/after relocation

```
0000000 <main>:  
    . . .  
    e: 83 ec 04      sub    $0x4,%esp  
    11: e8 fc ff ff ff call   12 <main+0x12>  
          12: R_386_PC32 swap  
    16: 83 c4 04      add    $0x4,%esp  
    . . .
```

Return address  
 $0x8048396 + 0x1a$   
= 0x80483b0  
Location of swap function

```
08048380 <main>:  
08048380: 8d 4c 24 04      lea    0x4(%esp),%ecx  
08048384: 83 e4 f0      and    $0xffffffff0,%esp  
08048387: ff 71 fc      pushl  0xfffffffffc(%ecx)  
0804838a: 55              push   %ebp  
0804838b: 89 e5          mov    %esp,%ebp  
0804838d: 51              push   %ecx  
0804838e: 83 ec 04      sub    $0x4,%esp  
08048391: e8 1a 00 00 00  call   80483b0 <swap>  
08048396: 83 c4 04      add    $0x4,%esp  
08048399: 31 c0          xor    %eax,%eax  
0804839b: 59              pop    %ecx  
0804839c: 5d              pop    %ebp  
0804839d: 8d 61 fc      lea    0xfffffffffc(%ecx),%esp  
080483a0: c3              ret
```

R\_386\_PC32  
relocate a  
reference that  
uses a 32-bit  
PC-relative  
address.

Effective address  
= PC + instruction  
encoded address

.text + .offset =  $0x8048380 + 0x12 = 0x8048392 \rightarrow$  reference address  
Call + 32-bit ref – reference address =  $0x80483b0 + -4 - 0x8048392 = 0x1a$